

## Solutions to problem set 3

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**Problem 1: Gram matrix**

Let  $\mathbf{A} \in \mathbb{C}^{M \times N}$  be the matrix whose  $k$ th column is  $\mathbf{a}_k$ . Since the entry in the  $k$ th row and in the  $l$ th column is  $\langle \mathbf{a}_k, \mathbf{a}_l \rangle = \mathbf{a}_l^H \mathbf{a}_k$ , the Gram matrix can be written as  $\mathbf{G} = \mathbf{A}^H \mathbf{A}$ . Therefore, it holds that

$$\mathbf{G}^H = (\mathbf{A}^H \mathbf{A})^H = \mathbf{A}^H (\mathbf{A}^H)^H = \mathbf{A}^H \mathbf{A} = \mathbf{G},$$

which shows that  $\mathbf{G}$  is Hermitian. Moreover, for all  $\mathbf{x} \in \mathbb{C}^M$ , we have

$$\mathbf{x}^H \mathbf{G} \mathbf{x} = \mathbf{x}^H \mathbf{A}^H \mathbf{A} \mathbf{x} = (\mathbf{A} \mathbf{x})^H \mathbf{A} \mathbf{x} = \|\mathbf{A} \mathbf{x}\|^2 \geq 0.$$

Therefore,  $\mathbf{G}$  is positive-semidefinite. Since  $\|\cdot\|$  is a norm,  $\mathbf{x}^H \mathbf{G} \mathbf{x} = 0$  implies that  $\mathbf{A} \mathbf{x} = 0$ . So, if  $\mathbf{A}$  has linearly independent columns,  $\mathbf{A} \mathbf{x} = 0$  implies that  $\mathbf{x} = 0$ , and  $\mathbf{G}$  is positive-definite.

**Problem 2: “ $l_0$ -norm”**

For a vector  $\mathbf{x} \in \mathbb{C}^N$ , we have that

$$\|2\mathbf{x}\|_0 = \|\mathbf{x}\|_0,$$

which shows that the homogeneity property of a norm is not satisfied. Therefore,  $\|\cdot\|_0$  is not a norm for  $\mathbb{C}^N$ . Its name and notation come from an abuse of terminology, taking  $p = 0$  in the definition of the  $l_p$ -norm  $\|\cdot\|_p$  which is commonly defined for  $p \in [1, \infty]$  as

$$\|\mathbf{x}\|_p^p = \sum_{k=1}^N |x_k|^p.$$

**Problem 3: Fat matrix inversion**

1. The equation  $\mathbf{A} \mathbf{x} = \mathbf{y}$  has infinitely many solutions if  $\mathbf{y}$  is in the column range space of  $\mathbf{A}$ . This is guaranteed when  $\text{rank} \mathbf{A} = S$ , i.e., when all rows of  $\mathbf{A}$  are linearly independent (or equivalently,  $S$  columns of  $\mathbf{A}$  are linearly independent). The equation  $\mathbf{A} \mathbf{x} = \mathbf{y}$  has no solution if  $\mathbf{y}$  is not in the column range space of  $\mathbf{A}$ . This can only happen when  $\text{rank} \mathbf{A} < S$ , i.e., when  $\mathbf{A}$  has linearly dependent rows (or equivalently, less than  $S$  linearly independent columns).
2. Solving the equation  $\mathbf{A} \mathbf{x} = \mathbf{y}$  under the constraint that  $x_j = 0$  for  $j \notin S$  amounts to solving the equation  $\tilde{\mathbf{A}} \tilde{\mathbf{x}} = \mathbf{y}$ , where  $\tilde{\mathbf{A}}$  is the  $S \times S$  matrix obtained by removing the  $N - S$  columns of  $\mathbf{A}$  that are indexed by  $S^c$  and where the unknown  $\tilde{\mathbf{x}}$  is an  $S$ -dimensional vector. The equation  $\tilde{\mathbf{A}} \tilde{\mathbf{x}} = \mathbf{y}$  has exactly one solution if  $\det \tilde{\mathbf{A}} \neq 0$ , i.e., if  $\tilde{\mathbf{A}}$  has full rank. Therefore, the equation  $\mathbf{A} \mathbf{x} = \mathbf{y}$  has exactly one solution if the columns  $\{\mathbf{a}_j\}_{j \in S}$  indexed by  $S$  are linearly independent.

#### Problem 4: Compressed sensing

Intuitively, we can write  $\mathbf{x}$  in terms of its discrete derivative. We can write  $\mathbf{x} = \mathbf{U}\mathbf{x}'$ , where

$$\mathbf{U} = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ 1 & 1 & 0 & \dots & 0 \\ 1 & 1 & 1 & \dots & 0 \\ & & & \ddots & \\ 1 & 1 & 1 & \dots & 1 \end{bmatrix} \in \mathbb{R}^{N \times N}$$

is the discrete integration matrix and  $\mathbf{x}' = [\underbrace{\alpha_1 0 \dots 0}_{\text{block1}} \underbrace{\alpha_2 - \alpha_1 0 \dots 0}_{\text{block2}} \dots \underbrace{\alpha_s - \alpha_{s-1} 0 \dots 0}_{\text{blocks}}]^\top \in \mathbb{R}^N$  is an  $s$ -sparse vector. We can then recover  $\mathbf{x}'$  from  $\mathbf{y} = \mathbf{A}\mathbf{x} = \mathbf{A}\mathbf{U}\mathbf{x}'$  by solving

$$\text{minimize}_{\hat{\mathbf{x}}} \|\hat{\mathbf{x}}\|_1 \text{ subject to } \mathbf{y} = \mathbf{A}\mathbf{U}\hat{\mathbf{x}}.$$

#### Problem 5: P0 recovery algorithm

See `pzero.ipynb`

#### Problem 6: Coherence in sines and spikes

By definition:

$$\mu(\mathbf{D}) = \max_{k,l} |\langle \mathbf{d}_k, \mathbf{d}_l \rangle|.$$

Observe that

$$\langle \mathbf{d}_k, \mathbf{d}_l \rangle = \begin{cases} 0, & k, l \leq M \text{ or } k, l > M \\ \frac{1}{\sqrt{M}} |e^{2\pi i k(l-M)/M}| = \frac{1}{\sqrt{M}}, & k \leq M \text{ and } l > M. \end{cases}$$

Therefore,  $\mu(\mathbf{D}) = \frac{1}{\sqrt{M}}$ .

From the lecture we know that if

$$s < \frac{1}{2} \left( 1 + \frac{1}{\mu(\mathbf{D})} \right) = \frac{1}{2} (1 + \sqrt{M})$$

successful recovery via P0 or basis pursuit is guaranteed for every  $s$ -sparse signal.

#### Problem 7: Coherence in super-resolution

Let's compute the absolute value of the inner product between the columns  $l_1$  and  $l_2$  of  $\mathbf{F}_{10}$ :

$$|\langle \mathbf{f}_{l_1}, \mathbf{f}_{l_2} \rangle| = \frac{1}{N} \left| \sum_{k=-f_c}^{f_c} e^{i2\pi k(l_1-l_2)/N} \right| = \frac{1}{N} \left| \sum_{k=0}^{M-1} e^{i2\pi k l / N} \right|$$

where  $l = l_1 - l_2$  and  $M = 2f_c + 1$ .

Using the formula for the sum of geometric series:

$$\sum_{k=0}^{M-1} e^{ikx} = \frac{1 - e^{iMx}}{1 - e^{ix}} = \frac{\sin(Mx/2)}{\sin(x/2)} e^{ix(M-1)/2}$$

we find:

$$|\langle \mathbf{f}_{l_1}, \mathbf{f}_{l_2} \rangle| = \frac{1}{N} \left| \frac{\sin(\pi l M/N)}{\sin(\pi l/N)} \right|.$$

Therefore,

$$\mu(\mathbf{F}_{\text{lo}}) = \max_{l_1, l_2} |\langle \mathbf{f}_{l_1}, \mathbf{f}_{l_2} \rangle| = \max_l \frac{1}{N} \left| \frac{\sin(\pi l M/N)}{\sin(\pi l/N)} \right| \geq \frac{1}{N} \left| \frac{\sin(\pi M/N)}{\sin(\pi/N)} \right|.$$

For  $M = N/2$  and large  $N$  we can approximate the right hand side of the equation above as

$$\frac{1}{N} \left| \frac{\sin(\pi M/N)}{\sin(\pi/N)} \right| \approx \frac{1}{N} \frac{\sin(\pi/2)}{\pi/N} = \frac{1}{\pi}.$$

From the lecture we know that if

$$s < \frac{1}{2} \left( 1 + \frac{1}{\mu(\mathbf{F}_{\text{lo}})} \right)$$

successful recovery via P0 or basis pursuit is guaranteed for every  $s$ -sparse signal. Since  $\mu(\mathbf{F}_{\text{lo}}) \geq \frac{1}{\pi}$  the bound we get for  $s$  is no better than

$$s < \frac{1}{2}(1 + \pi) \approx 2.$$

This is very pessimistic: for large  $N$  and  $M = N/2$  the coherence-based bound only guarantees the successful recovery of only 2 spikes. It turns out, as we will see in the lecture, a much more optimistic bound may be obtained for low-frequency Fourier measurements when the signal is nonnegative.

### Problem 8: Super-resolution experiment

See `superres.ipynb`